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DA: Datamover Architecture for
the Internet Small Computer System Interface (iSCSI)

Status of This Memo

This memo provides information for the Internet community. It does not specify an Internet standard of any kind. Distribution of this memo is unlimited.

Abstract

The Internet Small Computer System Interface (iSCSI) is a SCSI transport protocol that maps the SCSI family of application protocols onto TCP/IP. Datamover Architecture for iSCSI (DA) defines an abstract model in which the movement of data between iSCSI end nodes is logically separated from the rest of the iSCSI protocol in order to allow iSCSI to adapt to innovations available in new IP transports. While DA defines the architectural functions required of the class of Datamover protocols, it does not define any specific Datamover protocols. Each such Datamover protocol, defined in a separate document, provides a reliable transport for all iSCSI PDUs, but actually moves the data required for certain iSCSI PDUs without involving the remote iSCSI layer itself. This document begins with an introduction of a few new abstractions, defines a layered architecture for iSCSI and Datamover protocols, and then models the interactions within an iSCSI end node between the iSCSI layer and the Datamover layer that happen in order to transparently perform remote data movement within an IP fabric. It is intended that this definition will help map iSCSI to generic Remote Direct Memory Access (RDMA)-capable IP fabrics in the future comprising TCP, the Stream Control Transmission Protocol (SCTP), and possibly other underlying network transport layers, such as InfiniBand.

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1. Motivation

1.1. Intent

There are relatively new standard protocols that enable Remote Direct Memory Access (RDMA) and Remote Direct Data Placement (RDDP) technologies to work over IP fabrics. The principal value proposition of these technologies is that they enable one end node to place data in the final intended buffer on the remote end node, thus eliminating the need for a receive path data copy that moves the data to its final location. The data copy avoidance in turn eliminates unnecessary memory bandwidth consumption, substantially decreases the reassembly buffer size requirements, and preserves CPU cycles that would otherwise be spent in copying.

The iSCSI specification [RFC3720] defines a very detailed data transfer model that employs SCSI Data-In PDUs, SCSI Data-Out PDUs, and R2T PDUs, in addition to the SCSI Command and SCSI Response PDUs that respectively create and conclude the task context for the data transfer. In the traditional iSCSI model, the iSCSI protocol layer plays the central role in pacing the data transfer and carrying out the ensuing data transfer itself. An alternative architecture would be for iSCSI to delegate a large part of this data transfer role to a separate protocol layer exclusively designed to move data, which in turn is possibly aided by a data movement and placement technology such as RDMA.

If iSCSI were operating in such RDMA environments, iSCSI would be shielded from the low-level data transfer mechanics but would only be privy to the conclusion of the requested data transfer. Thus, there would be an effective "off-loading" of the work that an iSCSI protocol layer is expected to perform, compared to today's iSCSI end nodes. For such RDMA environments, it is highly desirable that there be a standard architecture to separate the data movement part of the iSCSI protocol definition from the rest of the iSCSI functionality. This architecture precisely defines what a Datamover layer is and also describes the model of interactions between the iSCSI layer and the Datamover layer (Section 6). In order to satisfy this need, this document presents a Datamover Architecture for iSCSI (DA) and summarizes a reasonable model for interactions between the iSCSI layer and the Datamover layer for each of the iSCSI PDUs that are defined in [RFC3720]. Note that while DA is motivated by the advent of RDMA over TCP/IP technology, the architecture is not dependent on RDMA in its design. DA is intended to be a generic architectural framework for allowing different types of Datamovers based on different types of RDMA and transport protocols. Adoption of this model will help iSCSI proliferate into more environments.

1.2. Interpretation of Requirements

This document introduces certain architectural abstractions and builds an abstract functional interface model between iSCSI and Datamover protocol layers based on those abstractions. This architectural style is motivated by the following desires:

- a) Provide guidance to Datamover protocol designers with respect to the functional boundary between iSCSI and the Datamover protocols. This guidance is critical since a significant part of the [RFC3720] protocol definition is left unchanged by DA architecture and the iSCSI notions from [RFC3720] (e.g., tasks, ITTs) are leveraged by the Datamover protocol.
- b) Aid existing iSCSI implementations to rapidly adapt to DA architecture, largely by leveraging the architectural abstractions into implementation constructs -- e.g., functions, APIs, modules.

However, note that DA architecture does not intend to impose any implementation specifics per se. When a DA architectural concept (e.g., Operational Primitive) is described as mandatory ("MUST") or recommended ("SHOULD") of a layer (iSCSI or Datamover) in this document, the intent is that an implementation respectively MUST or SHOULD produce the same protocol action as what the model describes. Specifically, no implementation compliance in terms of names, modules or API arguments etc. is implied by this Architecture by such use of [RFC2119] terms, only a functional compliance is sought.

2. Definitions and Acronyms

2.1. Definitions

I/O Buffer - A buffer that is used in a SCSI Read or Write operation so that SCSI data may be sent from or received by the buffer.

Datamover protocol - A Datamover protocol is a data transfer wire protocol for iSCSI that meets the requirements stated in Section 6.

Datamover layer - A Datamover layer is a protocol layer within an end node that implements the Datamover protocol.

Datamover-assisted - An iSCSI connection is said to be "Datamover-assisted" when a Datamover layer is enabled for moving control and data information on that iSCSI connection.

The key words "MUST", "MUST NOT", "REQUIRED", "SHALL", "SHALL NOT", "SHOULD", "SHOULD NOT", "RECOMMENDED", "MAY", and "OPTIONAL" in this document are to be interpreted as described in [RFC2119].

2.2. Acronyms

Acronym	Definition
DA	Datamover Architecture for iSCSI
DDP	Direct Data Placement Protocol
DI	Datamover Interface
IANA	Internet Assigned Numbers Authority
IETF	Internet Engineering Task Force
I/O	Input - Output
IP	Internet Protocol
iSCSI	Internet SCSI
iSER	iSCSI Extensions for RDMA
ITT	Initiator Task Tag
LO	Leading Only
MPA	Marker PDU Aligned Framing for TCP
PDU	Protocol Data Unit
RDDP	Remote Direct Data Placement
RDMA	Remote Direct Memory Access
R2T	Ready To Transfer
R2TSN	Ready To Transfer Sequence Number
RDMA	Remote Direct Memory Access
RDMAP	Remote Direct Memory Access Protocol
RFC	Request For Comments

SAM	SCSI Architecture Model
SCSI	Small Computer Systems Interface
SN	Sequence Number
SNACK	Selective Negative Acknowledgment - also Sequence Number Acknowledgement for Data
TCP	Transmission Control Protocol
TTT	Target Transfer Tag

3. Architectural Layering of iSCSI and Datamover Layers

Figure 1 illustrates an example of the architectural layering of iSCSI and Datamover layers, in conjunction with a TCP/IP implementation of RDMA/DDP ([DDP]) layers in an iSCSI end node. Note that RDMA/DDP/MPA and TCP protocol layers are shown here only as an example, and in reality, DA is completely oblivious to protocol layers below the Datamover layer. The RDMA/DDP/MPA protocol stack provides a generic transport service with direct data placement. There is no need to tailor the implementation of this protocol stack to the specific ULP to benefit from these services.

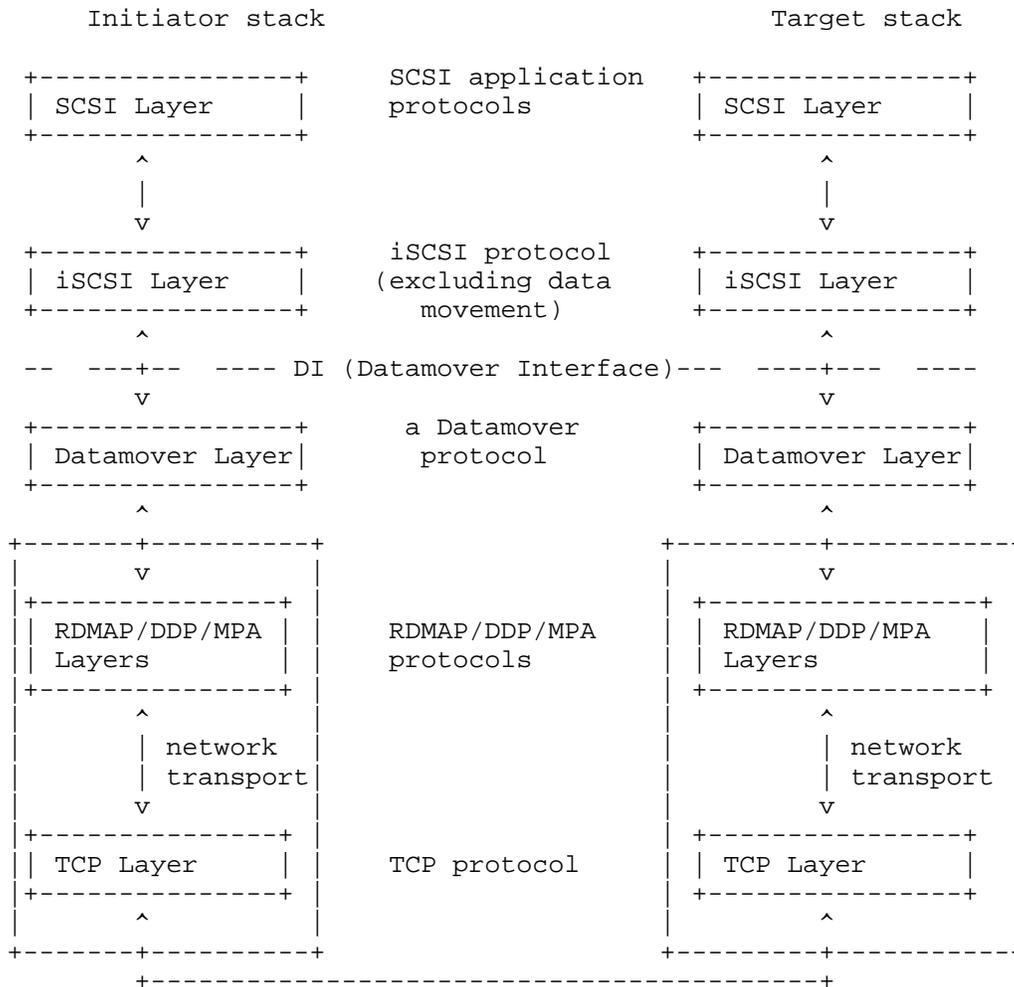


Figure 1. Datamover Architecture Diagram, with the RDMAP Example

The scope of this document is limited to:

1. Defining the notion of a Datamover layer and a Datamover protocol (Section 6).
2. Defining the functionality distribution between the iSCSI layer and the Datamover layer, along with the communication model between the two (Operational Primitives).

3. Modeling the interactions between the blocks labeled as "iSCSI Layer" and "Datamover Layer" in Figure 1 -- i.e., defining the interface labeled "DI" in the figure -- for each defined iSCSI PDU, based on the Operational Primitives.

4. Design Overview

This document discusses and defines a model for interactions between the iSCSI layer and a "Datamover layer" (see Section 6) operating within an iSCSI end node, presumably communicating with one or more iSCSI end nodes with similar layering. The model for interactions for handling different iSCSI operations is called the "Datamover Interface" (DI, Section 10), while the architecture itself is called the "Datamover Architecture for iSCSI" (DA). It is likely that the architecture will have implications on the Datamover wire protocols as DA places certain requirements and functionality expectations on the Datamover layer. However, this document itself neither defines any new wire protocol for the Datamover layer, nor any potential modifications to the iSCSI wire protocol to employ the Datamover layer. The scope of this document is strictly limited to specifying the architectural framework and the minimally required interactions that happen within an iSCSI end node to leverage the Datamover layer.

The design ideas behind DA can be summarized as follows:

- 1) DA defines an abstract functional interface model of the iSCSI layer's interactions with a Datamover layer below -- i.e., DA models the interactions between the logical "bottom" interface of iSCSI and the logical "top" interface of a Datamover.
- 2) DA guides the wire protocol for a Datamover layer by defining the iSCSI knowledge that the Datamover layer may utilize in its protocol definition (as an example, this document completely limits the notion of "iSCSI session" to the iSCSI layer).
- 3) DA is designed to allow implementation of the Datamover layer either in hardware or in software.
- 4) DA is not a wire protocol spec, but an architecture that also models the interactions between iSCSI and Datamover layers operating within an iSCSI end node.
- 5) DA by design seeks to model the iSCSI-Datamover interactions in a way that the modeling is independent of the specifics of either a particular iSCSI revision or an instantiation of a Datamover layer.

- 6) DA introduces and relies on the notion of a defined set of Operational Primitives (could be seen as entry point definitions in implementation terms) provided by each layer to the other to carry out the request-response interactions.
- 7) DA is intended to allow Datamover protocol definitions with minimal changes to existing iSCSI implementations.
- 8) DA is designed to allow the iSCSI layer to completely rely on the Datamover layer for all data transport needs.
- 9) DA models the architecturally required minimal interactions between an operational iSCSI layer and a Datamover layer to realize the iSCSI-transparent data movement. There may be several other interactions in a typical implementation in order to bootstrap a Datamover layer (or an iSCSI layer) into operation, but they are outside the scope of this document.

Note that in summary, DA is architected to support many different Datamover protocols operating under the iSCSI layer. One such example of a Datamover protocol is iSER [iSER].

5. Architectural Concepts

5.1. iSCSI PDU Types

This section defines the iSCSI PDU classification terminology, as defined and used in this document. Out of the set of legal iSCSI PDUs defined in [RFC3720], as we will see in Section 5.1.1, the iSCSI layer does not request a SCSI Data-Out PDU carrying solicited data for transmission across the Datamover Interface per this architecture. For this reason, the SCSI Data-Out PDU carrying solicited data is excluded in the iSCSI PDU classification we introduce in this section (for SCSI Data-Out PDUs for unsolicited Data, see Section 5.1.2). The rest of the legal iSCSI PDUs that may be exchanged across the Datamover Interface are defined to consist of two classes:

- 1) iSCSI data-type PDUs
- 2) iSCSI control-type PDUs

5.1.1. iSCSI Data-Type PDUs

An iSCSI data-type PDU is defined as an iSCSI PDU that causes data transfer, transparent to the remote iSCSI layer, to take place between the peer iSCSI nodes on a Full Feature Phase iSCSI connection. A data-type PDU, when requested for transmission by the

sender iSCSI layer, results in the associated data transfer without the participation of the remote iSCSI layer, i.e., the PDU itself is not delivered as-is to the remote iSCSI layer. The following iSCSI PDUs constitute the set of iSCSI data-type PDUs:

- 1) SCSI Data-In PDU
- 2) R2T PDU

In an iSCSI end node structured as an iSCSI layer and a Datamover layer as defined in this document, the solicitation for Data-Out (i.e., R2T PDU) is not delivered to the initiator iSCSI layer, per the definition of an iSCSI data-type PDU. The data transfer is instead performed via the mechanisms known to the Datamover layer (e.g., RDMA Read). This in turn implies that a SCSI Data-Out PDU for solicited data is never requested for transmission across the Datamover Interface at the initiator.

5.1.2. iSCSI Control-Type PDUs

Any iSCSI PDU that is not an iSCSI data-type PDU and also not a solicited SCSI Data-Out PDU is defined as an iSCSI control-type PDU. Specifically, note that SCSI Data-Out PDUs for unsolicited Data are defined as iSCSI control-type PDUs.

5.2. Data_Descriptor

A Data_Descriptor is an information element that describes an iSCSI/SCSI data buffer, provided by the iSCSI layer to its local Datamover layer or provided by the Datamover layer to its local iSCSI layer for identifying the data associated respectively with the requested or completed operation.

In implementation terms, a Data_Descriptor may be a scatter-gather list describing a local buffer, the exact structure of which is subject to the constraints imposed by the operating environment on the local iSCSI node.

5.3. Connection_Handle

A Connection_Handle is an information element that identifies the particular iSCSI connection for which an inbound or outbound iSCSI PDU is intended. A connection handle is unique for a given pair of an iSCSI layer instance and a Datamover layer instance. The Connection_Handle qualifier is used in all invocations of any Operational Primitive for connection identification.

Note that the `Connection_Handle` is conceptually different from the Connection Identifier (CID) defined by the iSCSI specification. While the CID is a unique identifier of an iSCSI connection within an iSCSI session, the uniqueness of the `Connection_Handle` extends to the entire iSCSI layer instance coupled with the Datamover layer instance, across possibly multiple iSCSI sessions.

In implementation terms, a `Connection_Handle` could be an opaque identifier exchanged between the iSCSI layer and the Datamover layer at the connection login time. One may also consider it to be similar in scope of uniqueness to a socket identifier. The exact structure and modalities of exchange of a `Connection_Handle` between the two layers is implementation-specific.

5.4. Operational Primitive

An Operational Primitive, in this document, is an abstract functional interface procedure that requests another layer perform a specific action on the requestor's behalf or notifies the other layer of some event. The Datamover Interface between an iSCSI layer instance and a Datamover layer instance within an iSCSI end node uses a set of Operational Primitives to define the functional interface between the two layers. Note that not every invocation of an Operational Primitive may elicit a response from the requested layer. This document describes the types of Operational Primitives that are implicitly required and provided by the iSCSI protocol layer as defined in [RFC3720], and the semantics of these Primitives.

Note that ownership of buffers and data structures is likely to be exchanged between the iSCSI layer and its local Datamover layer in invoking the Operational Primitives defined in this architecture. The buffer management details, including how buffers are allocated and released, are implementation-specific and thus are outside the scope of this document.

Each Operational Primitive invocation needs a certain "information context" (e.g., `Connection_Handle`) for performing the specific action being requested. The required information context is described in this document by a listing of "qualifiers" on each invocation, in the style of function call arguments. There is no specific implementation implied in this notation. The "qualifiers" of any Operational Primitive invocation specified in this document thus represent the mandatory information context that the Operational Primitive invocation MUST consider in performing the action. While the qualifiers are required, the method of realizing the qualifiers (passed synchronously with invocation, or retrieved from task context, or retrieved from shared memory etc.) is really up to the implementations.

When an Operational Primitive implementation is described as mandatory ("MUST") or recommended ("SHOULD") of a layer (iSCSI or Datamover) in this document, the intent is that an implementation respectively MUST or SHOULD produce the same protocol action as what the model describes.

5.5. Transport Connection

The term "Transport Connection" is used in this document as a generic term to represent the end-to-end logical connection as defined by the underlying reliable transport protocol. For this document, all instances of Transport Connection refer to a TCP connection.

6. Datamover Layer and Datamover Protocol

This section introduces the notion of a "Datamover layer" and "Datamover protocol" as meant in this document, and defines the requirements on a Datamover protocol.

A Datamover layer is the implementation component that realizes a Datamover protocol functionality in an iSCSI-capable end node in communicating with other iSCSI end nodes with similar capabilities. More specifically, a "Datamover layer" MUST provide the following functionality and the "Datamover protocol" MUST consist of the wire protocol required to realize the following functionality:

- 1) guarantee that all the necessary data transfers take place when the local iSCSI layer requests transmitting a command (in order to complete a SCSI command, for an initiator), or sending/receiving an iSCSI data sequence (in order to complete part of a SCSI command for a target).
- 2) transport an iSCSI control-type PDU as-is to the peer Datamover layer when requested to do so by the local iSCSI layer.
- 3) provide notification and delivery to the iSCSI layer upon arrival of an iSCSI control-type PDU.
- 4) provide an initiator-to-target data acknowledgement of SCSI read data back to the target iSCSI layer, when requested.
- 5) provide an asynchronous notification upon completion of a requested data transfer operation that moved data without involving the iSCSI layer.
- 6) place the SCSI data into the I/O buffers or pick up the SCSI data for transmission out of the data buffers that the iSCSI layer had requested to be used for a SCSI I/O.

- 7) provide an error-free (i.e., must have at least the same level of assurance of data integrity as the CRC32C iSCSI data digest), reliable, in-order delivery transport mechanism over IP networks in performing the data transfer, and asynchronously notify the iSCSI layer upon iSCSI connection termination.

Note that this architecture expects that each compliant Datamover protocol will define the precise means of satisfying the requirements specified in this section.

In order to meet the functional requirements listed in this section, certain Datamover protocols may require pre-posted buffers from the local iSCSI protocol layer via mechanisms outside the scope of this document. In some implementations, the absence of such buffers may result in a connection failure. Datamover protocols may also realize these functional requirements via methods not explicitly listed in this document.

7. Functional Overview

This section presents an overview of the functional interactions between the iSCSI layer and the Datamover layer as intended by this Architecture.

7.1. Startup

The iSCSI Login Phase on an iSCSI connection occurs as defined in [RFC3720]. The Architecture assumes that at the end of the Login Phase, both the initiator and target, if they had so decided, transition the connection to being Datamover-assisted. The precise means of how an iSCSI initiator and an iSCSI target agree on having the connection Datamover-assisted is defined by the Datamover protocol. The only architectural requirement is that all iSCSI interactions in the iSCSI Full Feature Phase MUST be Datamover-assisted subject to the prior agreement, meaning that the Datamover protocol is in the iSCSI-to-iSCSI communication path below the iSCSI layer on either side as shown in Figure 1. DA defines the `Enable_Datamover` Operational Primitive (Section 8.6) to bring about this transition to a Datamover-assisted connection.

The Architecture also assumes that the Datamover layer may require a certain number of opaque local resources for making a connection Datamover-assisted. DA thus defines the `Allocate_Connection_Resources` Operational Primitive (Section 8.4) to model this interaction. This Primitive is intended to be invoked on each side once the two sides decide (as previously noted) to have the connection be Datamover-assisted. The expected sequence of Primitive invocations is depicted in Figures 2 and 3 in Section 13.2. Figures

4, 5, and 6 illustrate how the Primitives may be employed to deal with various legal login outcomes.

7.2. Full Feature Phase

All iSCSI peer communication in the Full Feature Phase happens through the Datamover layers if the iSCSI connection is Datamover-assisted. The Architecture assumes that a Datamover layer may require a certain number of opaque local resources for each new iSCSI task. In the normal course of execution, these task-level resources in the Datamover layer are assumed to be transparently allocated on each task initiation and deallocated on the conclusion of each task as appropriate. In exception scenarios however -- scenarios that do not yield a SCSI Response for each task such as ABORT TASK operation -- the Architecture assumes that the Datamover layer needs to be notified of the individual task terminations to aid its task-level resource management. DA thus defines the `Deallocate_Task_Resources` Operational Primitive (Section 8.9) to model this task-resource management. In specifying the ITT qualifier for the `Deallocate_Task_Resources` Primitive, the Architecture further assumes that the Datamover layer tracks its opaque task-level local resources by the iSCSI ITT. DA also defines `Send_Control` (Section 8.1), `Put_Data` (Section 8.2), `Get_Data` (Section 8.3), `Data_Completion_Notify` (Section 9.3), `Data_ACK_Notify` (Section 9.4), and `Control_Notify` (Section 9.1) Operational Primitives to model the various Full Feature Phase interactions.

Figures 9, 10, and 11 in Section 13.2 show some Full Feature Phase interactions -- SCSI Write task, SCSI Read task, and a SCSI Read Data acknowledgement, respectively. Figure 12 in Section 13.2 illustrates how an ABORT TASK operation can be modeled leading to deterministic resource cleanup on the Datamover layer.

7.3. Wrap-up

Once an iSCSI connection becomes Datamover-assisted, the connection continues in that state until the end of the Full Feature Phase, i.e., the termination of the connection. The Architecture assumes that when a connection is normally logged out, the Datamover layer needs to be notified so that its connection-level opaque resources (see Section 7.1) may be freed up. DA thus defines a `Connection_Terminate` Operational Primitive (Section 8.7) to model this interaction. The Architecture further assumes that when a connection termination happens without iSCSI layer's involvement (e.g., TCP RST), the Datamover layer is capable of locally cleaning up its task-level and connection-level resources before notifying the iSCSI layer of the fact. DA thus defines the

Connection_Terminate_Notify Operational Primitive (Section 9.2) to model this interaction.

Figures 7 and 8 in Section 13.2 illustrate the interactions between the iSCSI and Datamover layers in normal and unexpected connection termination scenarios.

8. Operational Primitives Provided by the Datamover Layer

While the iSCSI specification itself does not have a notion of Operational Primitives, any iSCSI layer implementing the iSCSI specification functionally requires the following Operational Primitives from its Datamover layer. Thus, any Datamover protocol compliant with this architecture MUST implement the Operational Primitives described in this section. These Operational Primitives are invoked by the iSCSI layer as appropriate. Unless otherwise stated, all the following Operational Primitives may be used both on the initiator side and the target side. In general programming terminology, this set of Operational Primitives may be construed as "down calls".

- 1) Send_Control
- 2) Put_Data
- 3) Get_Data
- 4) Allocate_Connection_Resources
- 5) Deallocate_Connection_Resources
- 6) Enable_Datamover
- 7) Connection_Terminate
- 8) Notice_Key_Values
- 9) Deallocate_Task_Resources

8.1. Send_Control

Input qualifiers: Connection_Handle, iSCSI PDU-specific qualifiers

Return Results: Not specified.

An iSCSI layer requests that its local Datamover layer transmit an iSCSI control-type PDU to the peer iSCSI layer operating in the remote iSCSI node by this Operational Primitive. The Datamover layer

performs the requested operation, and may add its own protocol headers in doing so. The iSCSI layer MUST NOT invoke the Send_Control Operational Primitive on an iSCSI connection that is not yet Datamover-assisted.

An initiator iSCSI layer requesting the transfer of a SCSI Command PDU or a target iSCSI layer requesting the transfer of a SCSI response PDU are examples of invoking the Send_Control Operational Primitive. As Section 10.3.1 illustrates later on, the iSCSI PDU-specific qualifiers in this example are: BHS and AHS, DataDescriptorOut, DataDescriptorIn, ImmediateDataSize, and UnsolicitedDataSize.

8.2. Put_Data

Input qualifiers: Connection_Handle, contents of a SCSI Data-In PDU header, Data_Descriptor, Notify_Enable

Return Results: Not specified.

An iSCSI layer requests that its local Datamover layer transmit the data identified by the Data_Descriptor for the SCSI Data-In PDU to the peer iSCSI layer on the remote iSCSI node by this Operational Primitive. The Datamover layer performs the operation by using its own protocol means, completely transparent to the remote iSCSI layer. The iSCSI layer MUST NOT invoke the Put_Data Operational Primitive on an iSCSI connection that is not yet Datamover-assisted.

The Notify_Enable qualifier is used to request the local Datamover layer to generate or not generate the eventual local completion notification to the iSCSI layer for this Put_Data invocation. For detailed semantics of this qualifier, see Section 9.3.

A Put_Data Primitive may only be invoked by an iSCSI layer on the target to its local Datamover layer.

A target iSCSI layer requesting the transfer of an iSCSI read data sequence (also known as a read burst) is an example of invoking the Put_Data Operational Primitive.

8.3. Get_Data

Input qualifiers: Connection_Handle, contents of an R2T PDU, Data_Descriptor, Notify_Enable

Return Results: Not specified.

An iSCSI layer requests that its local Datamover layer retrieve certain data identified by the R2T PDU from the peer iSCSI layer on the remote iSCSI node and place it into the buffer identified by the Data_Descriptor by invoking this Operational Primitive. The Datamover layer performs the operation by using its own protocol means, completely transparent to the remote iSCSI layer. The iSCSI layer MUST NOT invoke the Get_Data Operational Primitive on an iSCSI connection that is not yet Datamover-assisted.

The Notify_Enable qualifier is used to request that the local Datamover layer generate or not generate the eventual local completion notification to the iSCSI layer for this Get_Data invocation. For detailed semantics of this qualifier, see Section 9.3.

A Get_Data Primitive may only be invoked by an iSCSI layer on the target to its local Datamover layer.

A target iSCSI layer requesting the transfer of an iSCSI write data sequence (also known as a write burst) is an example of invoking the Get_Data Operational Primitive.

8.4. Allocate_Connection_Resources

Input qualifiers: Connection_Handle[, Resource_Descriptor]

Return Results: Status.

By invoking this Operational Primitive, an iSCSI layer requests that its local Datamover layer perform all the Datamover-specific resource allocations required for the Full Feature Phase of an iSCSI connection. The Connection_Handle identifies the connection for which the iSCSI layer is requesting resources to be allocated. Allocation of these resources is a step towards eventually transitioning the connection to become a Datamover-assisted iSCSI connection. Note that the Datamover layer however does not allocate any Datamover-specific task-level resources upon invocation of this Primitive.

An iSCSI layer, in addition, optionally specifies the implementation-specific resource requirements for the iSCSI connection to the Datamover layer by passing an input qualifier called Resource_Descriptor. The exact structure of a Resource_Descriptor is implementation-dependent, and hence structurally opaque to DA.

A return result of Status=success means that the Allocate_Connection_Resources invocation corresponding to that

Connection_Handle succeeded. If an Allocate_Connection_Resources invocation is made for a Connection_Handle for which an earlier invocation succeeded, the return Status must be success and the request will be ignored by the Datamover layer. A return result of Status=failure means that the Allocate_Connection_Resources invocation corresponding to that Connection_Handle failed. There MUST NOT be more than one Allocate_Connection_Resources Primitive invocation outstanding for a given Connection_Handle at any time.

The iSCSI layer must invoke the Allocate_Connection_Resources Primitive before the invocation of the Enable_Datamover Primitive.

8.5. Deallocate_Connection_Resources

Input qualifiers: Connection_Handle

Return Results: Not specified.

By invoking this Operational Primitive, an iSCSI layer requests that its local Datamover layer deallocate all the Datamover-specific resources that may have been allocated earlier for the Transport Connection identified by the Connection_Handle. The iSCSI layer may invoke this Operational Primitive when the Datamover-specific resources associated with the Connection_Handle are no longer necessary (such as the Login failure of the corresponding iSCSI connection).

8.6. Enable_Datamover

Input qualifiers: Connection_Handle, Transport_Connection_Descriptor [, Final_Login_Response_PDU]

Return Results: Not specified.

By invoking this Operational Primitive, an iSCSI layer requests that its local Datamover layer assist all further iSCSI exchanges on the iSCSI connection (i.e., to make the connection Datamover-assisted) identified by the Connection_Handle, for which the Datamover-specific resource allocation was earlier made. The iSCSI layer MUST NOT invoke the Enable_Datamover Operational Primitive for an iSCSI connection unless there is a corresponding prior resource allocation.

The Final_Login_Response_PDU input qualifier is applicable only for a target, and contains the final Login Response that concludes the iSCSI Login Phase and which must be sent as a byte stream as expected by the initiator iSCSI layer. When this qualifier is used, the target-Datamover layer MUST transmit this final Login Response before Datamover assistance is enabled for the Transport Connection.

The iSCSI layer identifies the specific Transport Connection associated with the Connection_Handle to the Datamover layer by specifying the Transport_Connection_Descriptor. The exact structure of this Descriptor is implementation-dependent.

8.7. Connection_Terminate

Input qualifiers: Connection_Handle

Return Results: Not specified.

By invoking this Operational Primitive, an iSCSI layer requests that its local Datamover layer terminate the Transport Connection and deallocate all the connection and task resources associated with the Connection_Handle. When this Operational Primitive invocation returns to the iSCSI layer, the iSCSI layer may assume the full ownership of all the iSCSI-level resources, e.g., I/O Buffers, associated with the connection. This Operational Primitive may be invoked only with a valid Connection_Handle, and the Transport Connection associated with the Connection_Handle must already be Datamover-assisted.

8.8. Notice_Key_Values

Input qualifiers: Connection_Handle, Number of keys, a list of Key-Value pairs.

Return Results: Not specified.

By invoking this Operational Primitive, an iSCSI layer requests that its local Datamover layer take note of the negotiated values of the listed keys for the Transport Connection. This Operational Primitive may be invoked only with a valid Connection_Handle, and the Key-Value pairs MUST be the current values that were successfully agreed upon by the iSCSI peers for the connection. The Datamover layer may use the values of the keys to aid the Datamover operation as it deems appropriate. The specific keys to be passed as input qualifiers and the point(s) in time this Operational Primitive is invoked are implementation-dependent.

8.9. Deallocate_Task_Resources

Input qualifiers: Connection_Handle, ITT

Return Results: Not specified.

By invoking this Operational Primitive, an iSCSI layer requests that its local Datamover layer deallocate all Datamover-specific resources

that earlier may have been allocated for the task identified by the ITT qualifier. The iSCSI layer uses this Operational Primitive during exception processing when one or more active tasks are to be terminated without corresponding SCSI Response PDUs. This Primitive MUST be invoked for each active task terminated without a SCSI Response PDU. This Primitive MUST NOT be invoked by the iSCSI layer when a SCSI Response PDU normally concludes a task. When a SCSI Response PDU normally concludes a task (even if the SCSI Status was not a success), the Datamover layer is assumed to have automatically deallocated all Datamover-specific task resources for that task. Refer to Section 7.2 for a related discussion on the Architectural assumptions on the task-level Datamover resource management, especially with respect to when the resources are assumed to be allocated.

9. Operational Primitives Provided by the iSCSI Layer

While the iSCSI specification itself does not have a notion of Operational Primitives, any iSCSI layer implementing the iSCSI specification would have to provide the following Operational Primitives to its local Datamover layer. Thus, any iSCSI protocol implementation compliant with this architecture MUST implement the Operational Primitives described in this section. These Operational Primitives are invoked by the Datamover layer as appropriate and when the iSCSI connection is Datamover-assisted. Unless otherwise stated, all the following Operational Primitives may be used both on the initiator side and the target side. In general programming terminology, this set of Operational Primitives may be construed as "up calls".

- 1) Control_Notify
- 2) Connection_Terminate_Notify
- 3) Data_Completion_Notify
- 4) Data_ACK_Notify

9.1. Control_Notify

Input qualifiers: Connection_Handle, an iSCSI control-type PDU.

Return Results: Not specified.

A Datamover layer notifies its local iSCSI layer, via this Operational Primitive, of the arrival of an iSCSI control-type PDU from the peer Datamover layer on the remote iSCSI node. The iSCSI layer processes the control-type PDU as defined in [RFC3720].

A target iSCSI layer being notified of the arrival of a SCSI command is an example of invoking the Control_Notify Operational Primitive.

Note that implementations may choose to describe the "iSCSI control-type PDU" qualifier in this notification using a Data_Descriptor (Section 5.2) and not necessarily one contiguous buffer.

9.2. Connection_Terminate_Notify

Input qualifiers: Connection_Handle

Return Results: Not specified.

A Datamover layer notifies its local iSCSI layer on an unsolicited termination or failure of an iSCSI connection providing the Connection_Handle associated with the iSCSI Connection. The iSCSI layer MUST consider the Connection_Handle to be invalid upon being so notified. The iSCSI layer processes the connection termination as defined in [RFC3720]. The Datamover layer MUST deallocate the connection and task resources associated with the terminated connection before notifying the iSCSI layer of the termination via this Operational Primitive.

A target iSCSI layer is notified of an ungraceful connection termination by the Datamover layer when the underlying Transport Connection is torn down. Such a Connection_Terminate_Notify Operational Primitive may be triggered, for example, by a TCP RESET in cases where the underlying Transport Connection uses TCP.

9.3. Data_Completion_Notify

Input qualifiers: Connection_Handle, ITT, SN

Return Results: Not specified.

A Datamover layer notifies its local iSCSI layer on completing the retrieval of the data or upon sending the data, as requested in a prior iSCSI data-type PDU, from/to the peer Datamover layer on the remote iSCSI node via this Operational Primitive. The iSCSI layer processes the operation as defined in [RFC3720].

SN may be either the DataSN associated with the SCSI Data-In PDU or R2TSN associated with the R2T PDU depending on the SCSI operation. Note that, for targets, a TTT (see [RFC3720]) could have been specified instead of an SN. However, the considered choice was to leave the SN to be the qualifier for two reasons -- a) it is generic and applicable to initiators and targets as well as Data-In and Data-Out, and b) having both SN and TTT qualifiers for the

notification is considered onerous on the Datamover layer, in terms of state maintenance for each completion notification. The implication of this choice is that iSCSI target implementations will have to adapt to using the ITT-SN tuple in associating the solicited data to the appropriate task, rather than the ITT-TTT tuple for doing the same.

If `Notify_Enable` is set in either a `Put_Data` or a `Get_Data` invocation, the Datamover layer MUST invoke the `Data_Completion_Notify` Operational Primitive upon completing that requested data transfer. If the `Notify_Enable` was cleared in either a `Put_Data` or a `Get_Data` invocation, the Datamover layer MUST NOT invoke the `Data_Completion_Notify` Operational Primitive upon completing that requested data transfer.

A `Data_Completion_Notify` invocation serves to notify the iSCSI layer of the `Put_Data` or `Get_Data` completion, respectively. As earlier noted in Sections 8.2 and 8.3, specific Datamover protocol definitions may restrict the usage scope of `Put_Data` and `Get_Data`, and thus implicitly the usage scope of `Data_Completion_Notify`.

A target iSCSI layer being notified of the retrieval of a write data sequence is an example of invoking the `Data_Completion_Notify` Operational Primitive.

9.4. `Data_ACK_Notify`

Input qualifiers: `Connection_Handle`, `ITT`, `DataSN`

Return Results: Not specified.

A target Datamover layer notifies its local iSCSI layer of the arrival of a previously requested data acknowledgement from the peer Datamover layer on the remote (initiator) iSCSI node via this Operational Primitive. The iSCSI layer processes the data acknowledgement notification as defined in [RFC3720].

A target iSCSI layer being notified of the arrival of a data acknowledgement for a certain SCSI Read data PDU is the only example of invoking the `Data_ACK_Notify` Operational Primitive.

10. Datamover Interface (DI)

10.1. Overview

This section describes the model of interactions between iSCSI and Datamover layers when the iSCSI connection is Datamover-assisted so the iSCSI layer may carry out the following:

- send iSCSI data-type PDUs and exchange iSCSI control-type PDUs, and
- handle asynchronous notifications such as completion of data sequence transfer and connection failure.

This chapter relies on the notion of Operational Primitives (Section 5.4) to define DI.

10.2. Interactions for Handling Asynchronous Notifications

10.2.1. Connection Termination

As stated in Section 9.2, the Datamover layer notifies the iSCSI layer of a failed or terminated connection via the `Connection_Terminate_Notify` Operational Primitive. The iSCSI layer MUST consider the connection unusable upon the invocation of this Primitive and handle the connection termination as specified in [RFC3720].

10.2.2. Data Transfer Completion

As stated in Section 9.3, the Datamover layer notifies the iSCSI layer of a completed data transfer operation via the `Data_Completion_Notify` Operational Primitive. The iSCSI layer processes the transfer completion as specified in [RFC3720].

10.2.2.1. Completion of a Requested SCSI Data Transfer

To notify the iSCSI layer of the completion of a requested iSCSI data-type PDU transfer, the Datamover layer uses the `Data_Completion_Notify` Operational Primitive with the following input qualifiers.

- a) `Connection_Handle`.
- b) `ITT`: Initiator Task Tag semantics as defined in [RFC3720].
- c) `SN`: `DataSN` for a SCSI Data-in/Data-out PDU, and `R2TSN` for an iSCSI R2T PDU. The semantics for both types of sequence numbers are as defined in [RFC3720].

The rationale for choosing `SN` is explained in Section 9.3.

Every invocation of the `Data_Completion_Notify` Operational Primitive MUST be preceded by an invocation of the `Put_Data` or `Get_Data` Operational Primitive with the `Notify_Enable` qualifier set by the iSCSI layer at an earlier point in time.

10.2.3. Data Acknowledgement

[RFC3720] allows the iSCSI targets to optionally solicit data acknowledgement from the initiator for one or more Data-In PDUs, via setting of the A-bit on a Data-In PDU. The Data_ACK_Notify Operational Primitive with the following input qualifiers is used by the target Datamover layer to notify the local iSCSI layer of the arrival of data acknowledgement of a previously solicited iSCSI read data acknowledgement. This Operational Primitive thus is applicable only to iSCSI targets.

- a) Connection_Handle.
- b) ITT: Initiator Task Tag semantics as defined in [RFC3720].
- c) DataSN: of the next SCSI Data-In PDU, which immediately follows the SCSI Data-In PDU with the A-bit set to which this notification corresponds, with semantics as defined in [RFC3720].

Every invocation of the Data_ACK_Notify Operational Primitive MUST be preceded by an invocation of the Put_Data Operational Primitive by the iSCSI target layer with the A-bit set to 1 at an earlier point in time.

10.3. Interactions for Sending an iSCSI PDU

This section discusses the model of interactions for sending each of the iSCSI PDUs defined in [RFC3720]. A Connection_Handle (see Section 5.3) is assumed to qualify each of these interactions so that the Datamover layer can route it to the appropriate Transport Connection. The qualifying Connection_Handle is not explicitly listed in the subsequent sections.

Note that the defined list of input qualifiers represents the semantically required set for the Datamover layer to consider in implementing the Primitive in each interaction described in this section (see Section 5.4 for an elaboration). Implementations may choose to deduce the qualifiers in ways that are optimized for the implementation specifics. Two examples of this are:

1. For SCSI command (Section 10.3.1), deducing the ImmediateDataSize input qualifier from the DataSegmentLength field of the SCSI Command PDU.
2. For SCSI Data-Out (Section 10.3.5.1), deducing the DataDescriptorOut input qualifier from the associated SCSI command invocation qualifiers (assuming such state is

maintained) in conjunction with BHS fields of the SCSI Data-Out PDU.

10.3.1. SCSI Command

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a SCSI Command PDU.

- a) BHS and AHS, if any, of the SCSI Command PDU as defined in [RFC3720].
- b) DataDescriptorOut: that defines the I/O Buffer meant for Data-Out for the entire command, in the case of a write or bidirectional command.
- c) DataDescriptorIn: that defines the I/O Buffer meant for Data-In for the entire command, in the case of a read or bidirectional command.
- d) ImmediateDataSize: that defines the number of octets of immediate unsolicited data for a write/bidirectional command.
- e) UnsolicitedDataSize: that defines the number of octets of immediate and non-immediate unsolicited data for a write/bidirectional command.

10.3.2. SCSI Response

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a SCSI Response PDU.

- a) BHS of the SCSI Response PDU as defined in [RFC3720].
- b) DataDescriptorStatus: that defines the iSCSI buffer that contains the sense and response information for the command.

10.3.3. Task Management Function Request

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a Task Management Function Request PDU.

- a) BHS of the Task Management Function Request PDU as defined in [RFC3720].

- b) DataDescriptorOut: that defines the I/O Buffer meant for Data-Out for the entire command, in the case of a write or bidirectional command. (Only valid if Function="TASK REASSIGN" - [RFC3720].)
- c) DataDescriptorIn: that defines the I/O Buffer meant for Data-In for the entire command, in the case of a read or bidirectional command. (Only valid if Function="TASK REASSIGN" - [RFC3720].)

10.3.4. Task Management Function Response

The Send_Control Operational Primitive with the following input qualifier is used for requesting the transmission of a Task Management Function Response PDU.

- a) BHS of the Task Management Function Response PDU as defined in [RFC3720].

10.3.5. SCSI Data-Out and SCSI Data-In

10.3.5.1. SCSI Data-Out

The Send_Control Operational Primitive with the following input qualifiers is used by the initiator iSCSI layer for requesting the transmission of a SCSI Data-Out PDU carrying the non-immediate unsolicited data.

- a) BHS of the SCSI Data-Out PDU as defined in [RFC3720].
- b) DataDescriptorOut: that defines the I/O Buffer with the Data-Out to be carried in the iSCSI data segment of the PDU.

10.3.5.2. SCSI Data-In

The Put_Data Operational Primitive with the following input qualifiers is used by the target iSCSI layer for requesting the transmission of the data carried by a SCSI Data-In PDU.

- a) BHS of the SCSI Data-In PDU as defined in [RFC3720].
- b) DataDescriptorIn: that defines the I/O Buffer with the Data-In being requested for transmission.

10.3.6. Ready To Transfer (R2T)

The Get_Data Operational Primitive with the following input qualifiers is used by the target iSCSI layer for requesting the retrieval of the data as specified by the semantic content of an R2T PDU.

- a) BHS of the Ready To Transfer PDU as defined in [RFC3720].
- b) DataDescriptorOut: that defines the I/O Buffer for the Data-Out being requested for retrieval.

10.3.7. Asynchronous Message

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of an Asynchronous Message PDU.

- a) BHS of the Asynchronous Message PDU as defined in [RFC3720].
- b) DataDescriptorSense: that defines an iSCSI buffer that contains the sense and iSCSI Event information.

10.3.8. Text Request

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a Text Request PDU.

- a) BHS of the Text Request PDU as defined in [RFC3720].
- b) DataDescriptorTextOut: that defines the iSCSI Text Request buffer.

10.3.9. Text Response

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a Text Response PDU.

- a) BHS of the Text Response PDU as defined in [RFC3720].
- b) DataDescriptorTextIn: that defines the iSCSI Text Response buffer.

10.3.10. Login Request

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a Login Request PDU.

- a) BHS of the Login Request PDU as defined in [RFC3720].
- b) DataDescriptorLoginRequest: that defines the iSCSI Login Request buffer.

Note that specific Datamover protocols may choose to disallow the standard DA Primitives from being used for the iSCSI Login Phase. When used in conjunction with such Datamover protocols, an attempt to send a Login Request via the Send_Control Operational Primitive invocation is clearly an error scenario, as the Login Request PDU is being sent while the connection is in the iSCSI Full Feature Phase. It is outside the scope of this document to specify the resulting implementation behavior in this case -- [RFC3720] already defines the error handling for this error scenario.

10.3.11. Login Response

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a Login Response PDU.

- a) BHS of the Login Response PDU as defined in [RFC3720].
- b) DataDescriptorLoginResponse: that defines the iSCSI Login Response buffer.

Note that specific Datamover protocols may choose to disallow the standard DA Primitives from being used for the iSCSI Login Phase. When used in conjunction with such Datamover protocols, an attempt to send a Login Response via the Send_Control Operational Primitive invocation is clearly an error scenario, as the Login Response PDU is being sent while in the iSCSI Full Feature Phase. It is outside the scope of this document to specify the resulting implementation behavior in this case -- [RFC3720] already defines the error handling for this error scenario.

10.3.12. Logout Command

The Send_Control Operational Primitive with the following input qualifier is used for requesting the transmission of a Logout Command PDU.

- a) BHS of the Logout Command PDU as defined in [RFC3720].

10.3.13. Logout Response

The Send_Control Operational Primitive with the following input qualifier is used for requesting the transmission of a Logout Response PDU.

- a) BHS of the Logout Response PDU as defined in [RFC3720].

10.3.14. SNACK Request

The Send_Control Operational Primitive with the following input qualifier is used for requesting the transmission of a SNACK Request PDU.

- a) BHS of the SNACK Request PDU as defined in [RFC3720].

10.3.15. Reject

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a Reject PDU.

- a) BHS of the Reject PDU as defined in [RFC3720].
- b) DataDescriptorReject: that defines the iSCSI Reject buffer.

10.3.16. NOP-Out

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a NOP-Out PDU.

- a) BHS of the NOP-Out PDU as defined in [RFC3720].
- b) DataDescriptorNOPOut: that defines the iSCSI Ping data buffer.

10.3.17. NOP-In

The Send_Control Operational Primitive with the following input qualifiers is used for requesting the transmission of a NOP-In PDU.

- a) BHS of the NOP-In PDU as defined in [RFC3720].
- b) DataDescriptorNOPIn: that defines the iSCSI Return Ping data buffer.

10.4. Interactions for Receiving an iSCSI PDU

The only PDUs that are received by an iSCSI layer operating on a Datamover layer are the iSCSI control-type PDUs. The Datamover layer delivers the iSCSI control-type PDUs as they arrive, qualifying each with the `Connection_Handle` (see Section 5.3) that identifies the iSCSI connection for which the PDU is meant. The subsequent processing of the iSCSI control-type PDUs proceeds as defined in [RFC3720].

10.4.1. General Control-Type PDU Notification

This sub-section describes the general mechanics applicable to several control-type PDUs. The following sub-sections note additional considerations for control-type PDUs that are not covered in this sub-section.

The `Control_Notify` Operational Primitive is used to notify the iSCSI layer of the arrival of the following iSCSI control-type PDUs: SCSI Command, SCSI Response, Task Management Function Request, Task Management Function Response, Asynchronous Message, Text Request, Text Response, Logout Command, Logout Response, SNACK, Reject, NOP-Out, NOP-In.

10.4.2. SCSI Data Transfer PDUs

10.4.2.1. SCSI Data-Out

The `Control_Notify` Operational Primitive is used to notify the iSCSI layer of the arrival of a SCSI Data-Out PDU carrying the non-immediate unsolicited data. Note however that the solicited SCSI Data-Out arriving on the target does not cause a notification to the iSCSI layer using the `Control_Notify` Primitive because the solicited SCSI Data-Out was not sent by the initiator iSCSI layer as control-type PDUs.

10.4.2.2. SCSI Data-In

The Datamover layer does not notify the iSCSI layer of the arrival of the SCSI Data-in at the initiator, because SCSI Data-in is an iSCSI data-type PDU (see section 5.1). The iSCSI layer at the initiator however may infer the arrival of the SCSI Data-In when it receives a subsequent notification of the SCSI Response PDU via a `Control_Notify` invocation.

While this document does not contemplate the possibility of a Data-In PDU being received at the initiator iSCSI layer, specific Datamover protocols may define how to deal with an unexpected inbound SCSI

Data-In PDU that may result in the initiator iSCSI layer receiving the Data-In PDU. This document leaves the details of handling this error scenario to the specific Datamover protocols, so each may define the appropriate error handling specific to the Datamover environment.

10.4.2.3. Ready To Transfer (R2T)

Because an R2T PDU is an iSCSI data-type PDU (see Section 5.1) that is not delivered as-is to the initiator iSCSI layer, the Datamover layer does not notify the iSCSI layer of the arrival of an R2T PDU. When an iSCSI node sends an R2T PDU to its local Datamover layer, the local and remote Datamover layers transparently bring about the data transfer requested by the R2T PDU.

While this document does not contemplate the possibility of an R2T PDU being received at the initiator iSCSI layer, specific Datamover protocols may define how to deal with an unexpected inbound R2T PDU that may result in the initiator iSCSI layer receiving the R2T PDU. This document leaves the details of handling this error scenario to the specific Datamover protocols, so each may define the appropriate error handling specific to the Datamover environment.

10.4.3. Login Request

The Control_Notify Operational Primitive is used for notifying the target iSCSI layer of the arrival of a Login Request PDU. Note that specific Datamover protocols may choose to disallow the standard DA Primitives from being used for the iSCSI Login Phase. When used in conjunction with such Datamover protocols, the arrival of a Login Request necessitating the Control_Notify Operational Primitive invocation is clearly an error scenario, as the Login Request PDU is arriving in the iSCSI Full Feature Phase. It is outside the scope of this document to specify the resulting implementation behavior in this case -- [RFC3720] already defines the error handling in this error scenario.

10.4.4. Login Response

The Control_Notify Operational Primitive is used to notify the initiator iSCSI layer of the arrival of a Login Response PDU. Note that specific Datamover protocols may choose to disallow the standard DA Primitives from being used for the iSCSI Login Phase. When used in conjunction with such Datamover protocols, the arrival of a Login Response necessitating the Control_Notify Operational Primitive invocation is clearly an error scenario, as the Login Response PDU is arriving in the iSCSI Full Feature Phase. It is outside the scope of this document to specify the resulting implementation behavior in

this case -- [RFC3720] already defines the error handling in this error scenario.

11. Security Considerations

11.1. Architectural Considerations

DA enables compliant iSCSI implementations to realize a control and data separation in the way they interact with their Datamover protocols. Note however that this separation does not imply a separation in transport mediums between control traffic and data traffic -- the basic iSCSI architecture with respect to tasks and PDU relationships to tasks remains unchanged. [RFC3720] defines several MUST requirements on ordering relationships across control and data for a given task besides a mandatory deterministic task allegiance model -- DA does not change this basic architecture (DA has a normative reference to [RFC3720]) for allow any additional flexibility in compliance in this area. To summarize, sending bulk data transfers (prompted by Put_Data and Get_Data Primitive invocations) on a different transport medium would be as ill-advised as sending just the Data-Out/Data-In PDUs on a different TCP connection in RFC 3720-based iSCSI implementations. Consequently, all the iSCSI-related security text in [RFC3723] is directly applicable to a DA-enabled iSCSI implementation.

Another area with security implications is the Datamover connection resource management model, which DA defines -- particularly the Allocate_Connection_Resources Primitive. An inadvertent realization of this model could leave an iSCSI implementation exposed to denial-of-service attacks. As Figures 2 and 3 in Section 13.2 illustrate, the most effective countermeasure to this potential attack consists of performing the Datamover resource allocation when the iSCSI layer is sufficiently far along in the iSCSI Login Phase that it is reasonably certain that the peer side is not an attacker. In particular, if the Login Phase includes a SecurityNegotiation stage, an iSCSI end node MUST defer the Datamover connection resource allocation (i.e., invoking the Allocate_Connection_Resources Primitive) to the LoginOperationalNegotiation stage [RFC3720] so that the resource allocation happens post-authentication. This considerably minimizes the potential for a denial-of service attack.

11.2. Wire Protocol Considerations

In view of the fact that the DA architecture itself does not define any new wire protocol or propose modifications to the existing protocols, there are no additional wire protocol security considerations in employing DA itself. However, a DA-compliant iSCSI implementation MUST comply with all the iSCSI-related requirements

stipulated in [RFC3723] and [RFC3720]. Note further that in realizing DA, each Datamover protocol must define and elaborate as appropriate on any additional security considerations resulting from the use of that Datamover protocol.

All Datamover protocol designers are strongly recommended to refer to [RDDPSEC] for the types of security issues to consider. While [RDDPSEC] elaborates on the security considerations applicable to an RDDP-based Datamover [iSER], the document is representative of the type of analysis of resource exhaustion and the application of countermeasures that need to be done for any Datamover protocol.

12. References

12.1. Normative References

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- [RFC3723] Aboba, B., Tseng, J., Walker, J., Rangan, V., and F. Travostino, "Securing Block Storage Protocols over IP", RFC 3723, April 2004.
- [RFC2119] Bradner, S., "Key words for use in RFCs to Indicate Requirement Levels", BCP 14, RFC 2119, March 1997.

12.2. Informative References

- [DDP] Shah, H., Pinkerton, J., Recio, R., and P. Culley, "Direct Data Placement over Reliable Transports", RFC 5041, October 2007.
- [iSER] Ko, M., Chadalapaka, M., Hufferd, J., Elzur, U., Shah, H., and P. Thaler, "Internet Small Computer System Interface (iSCSI) Extensions for Remote Direct Memory Access (RDMA)", RFC 5046, October 2007.
- [RDDPSEC] Pinkerton, J. and E. Deleganes, "Direct Data Placement Protocol (DDP) / Remote Direct Memory Access Protocol (RDMAP) Security", RFC 5042, October 2007.

Appendix A. Design Considerations and Examples

A.1. Design Considerations for a Datamover Protocol

This section discusses the specific considerations for RDMA-based and RDDP-based Datamover protocols.

- a) Note that the modeling of interactions for SCSI Data-Out (Section 10.3.5.1) is only used for unsolicited data transfer.
- b) The modeling of interactions for SNACK (Sections 10.3.14 and 10.4.1) is not expected to be used given that one of the design requirements on the Datamover is that it "guarantees an error-free, reliable, in-order transport mechanism" (Section 6). The interactions for sending and receiving a SNACK are nevertheless modeled in this document because the receiving iSCSI layer can deterministically deal with an inadvertent SNACK. This also shows the DA designers' intent that DI is not meant to filter certain types of PDUs.
- c) The onus is on a reliable Datamover (per requirements stated in Section 6) to realize end-to-end data acknowledgements via Datamover-specific means. In view of this, even use of data-ACK-type SNACKs are unnecessary. Consequently, an initiator may never request sending a SNACK Request in this model assuming that the proactive (timeout-driven) SNACK functionality is turned off in the legacy iSCSI code.
- d) Note that the current DA model for bootstrapping a `Connection_Handle` into service -- i.e., associating a new iSCSI connection with a `Connection_Handle` -- clearly implies that the iSCSI connection must already be in Full Feature Phase when the Datamover layer comes into the stack. This further implies that the iSCSI Login Phase must be carried out in the traditional "Byte streaming mode" with no assistance or involvement from the Datamover layer.

A.2. Examples of Datamover Interactions

The figures described in this section provide some examples of the usage of Operational Primitives in interactions between the iSCSI layer and the Datamover layer. The following abbreviations are used in this section.

Avail - Available

Abted - Aborted

Buf - I/O Buffer
Cmd - Command
Compl - Complete
Conn - Connection
Ctrl_Ntfy - Control_Notify
Dal_Tk_Res - Deallocate_Task_Resources
Data_Cmp_Nfy - Data_Completion_Notify
Data_ACK_Nfy - Data_ACK_Notify
DM - Datamover
Imm - Immediate
Snd_Ctrl - Send_Control
Msg - Message
Resp - Response
Sol - Solicited
TMF Req - Task Management Function Request
TMF Res - Task Management Function Response
Trans - Transfer
Unsol - Unsolicited

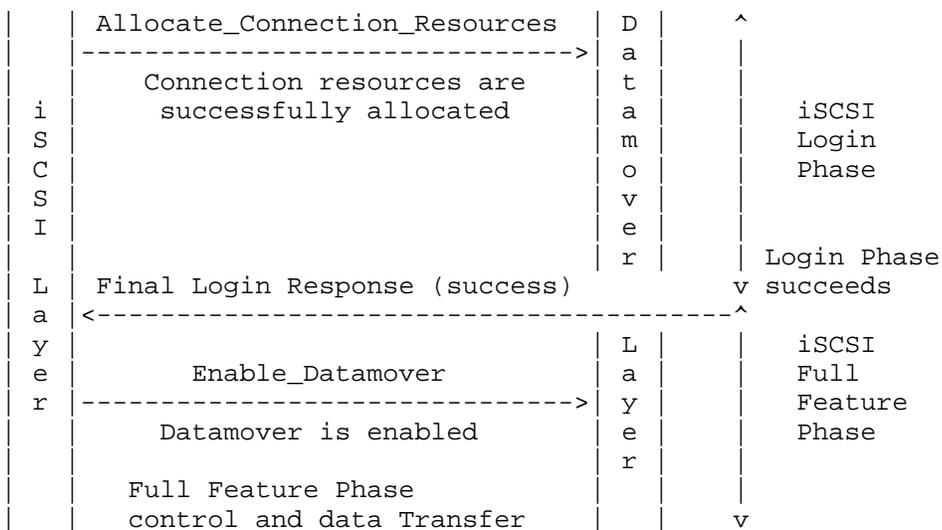


Figure 2. A Successful iSCSI Login on Initiator

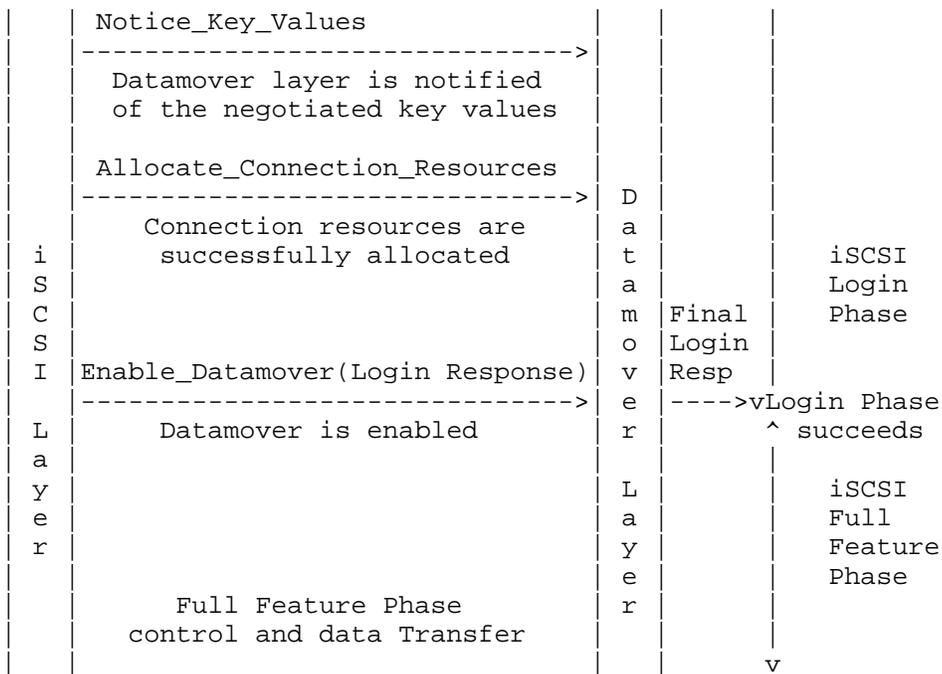


Figure 3. A Successful iSCSI Login on Target

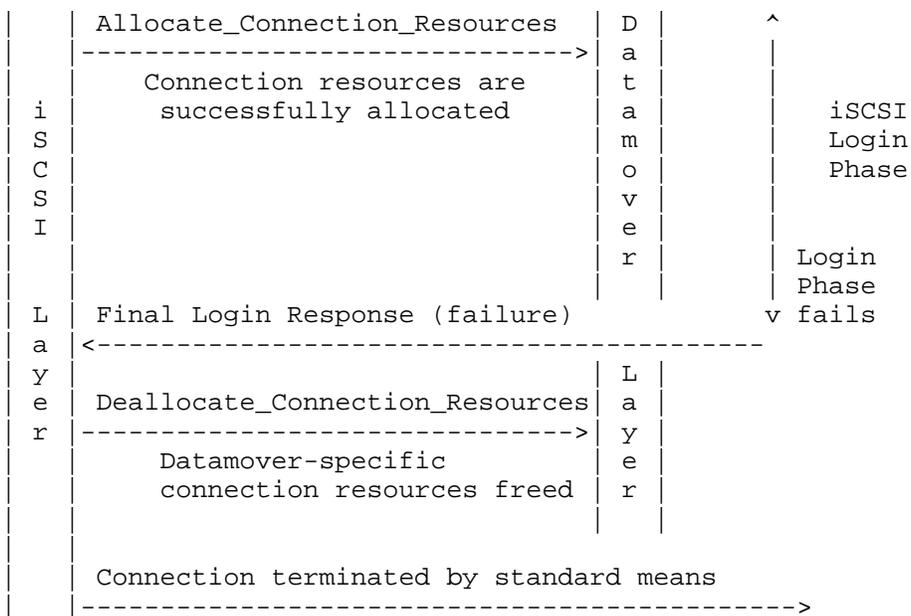


Figure 4. A Failed iSCSI Login on Initiator

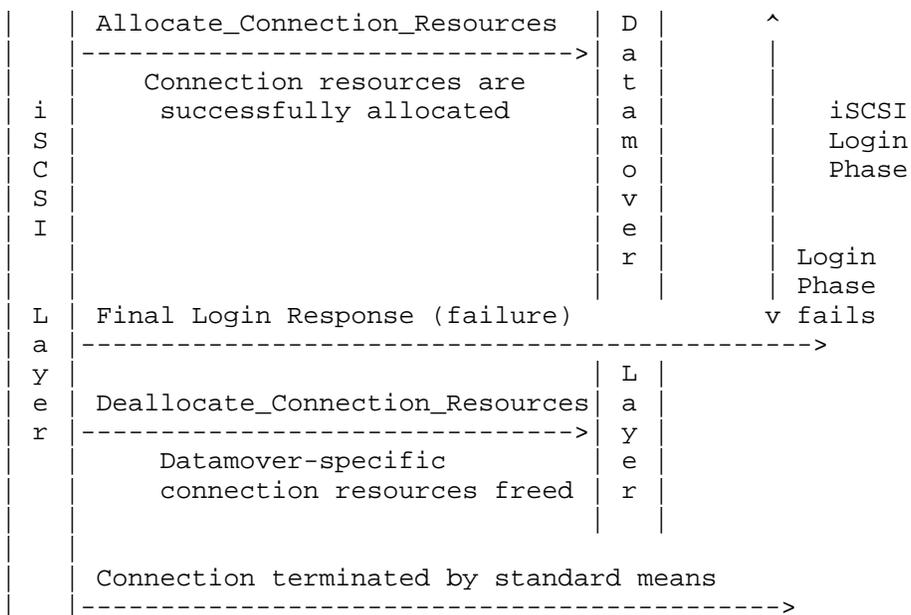


Figure 5. A Failed iSCSI Login on Target

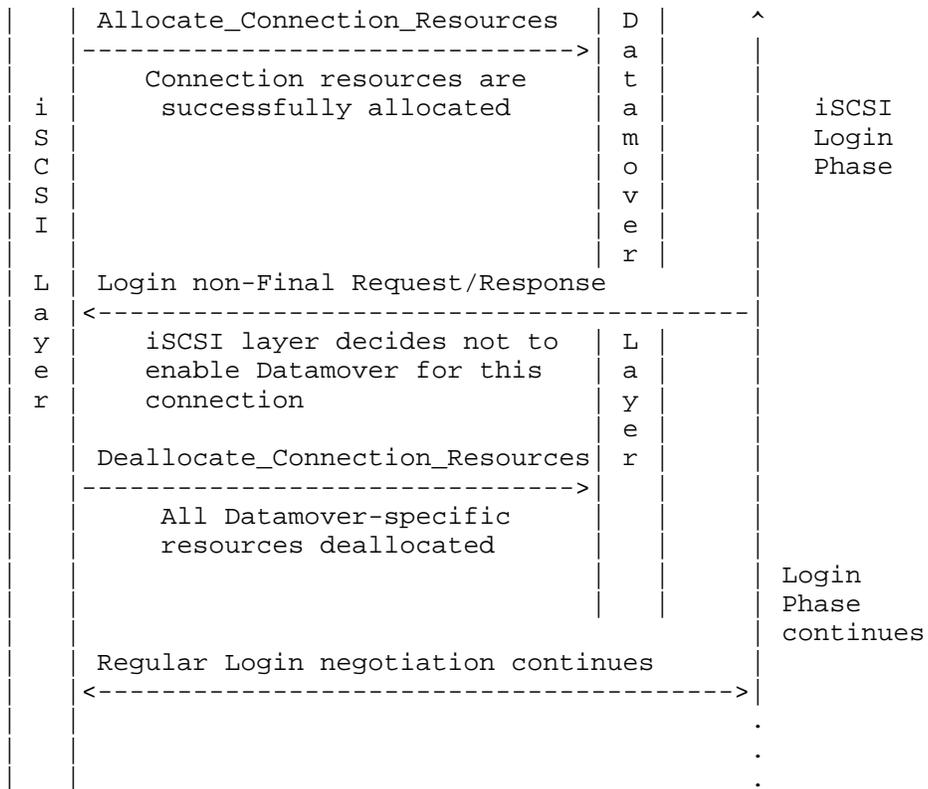


Figure 6. iSCSI Does Not Enable the Datamover

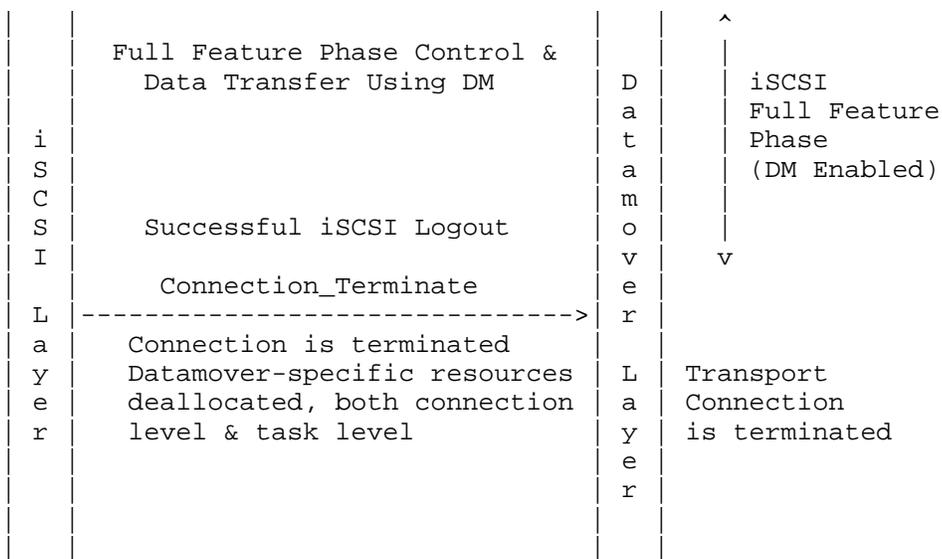


Figure 7. A Normal iSCSI Connection Termination

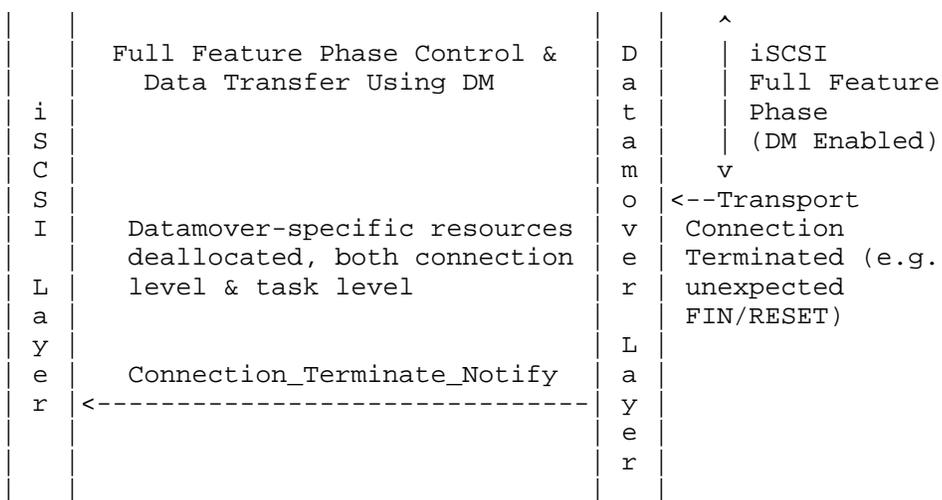


Figure 8. An Abnormal iSCSI Connection Termination

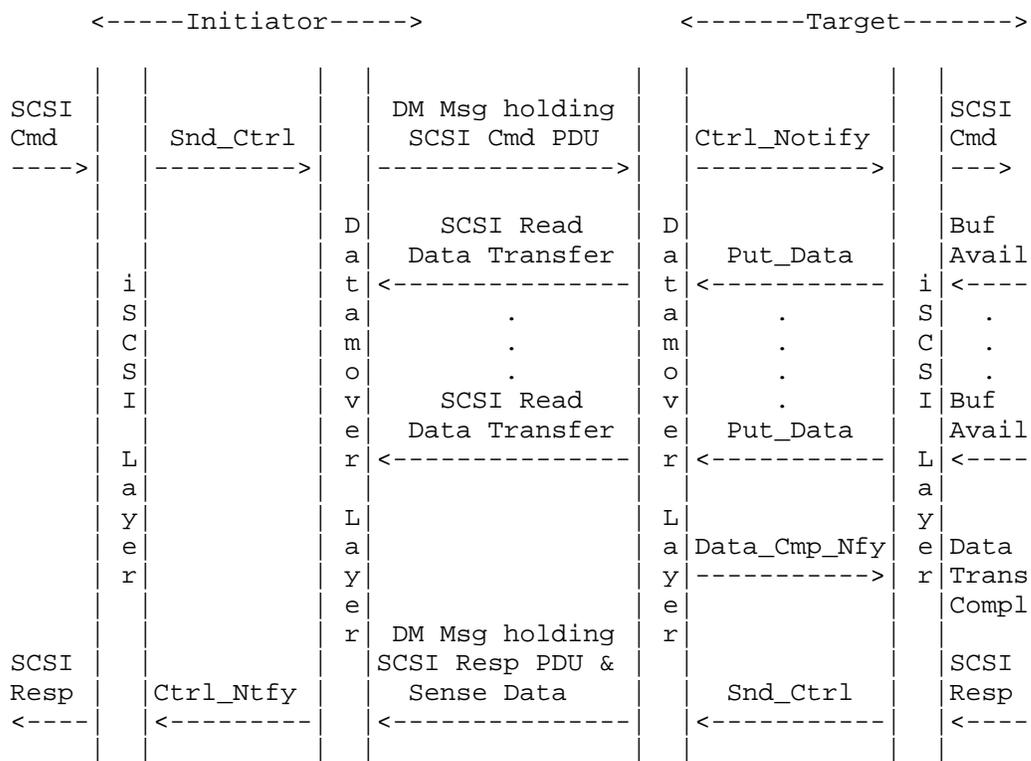


Figure 10. A SCSI Read Data Transfer

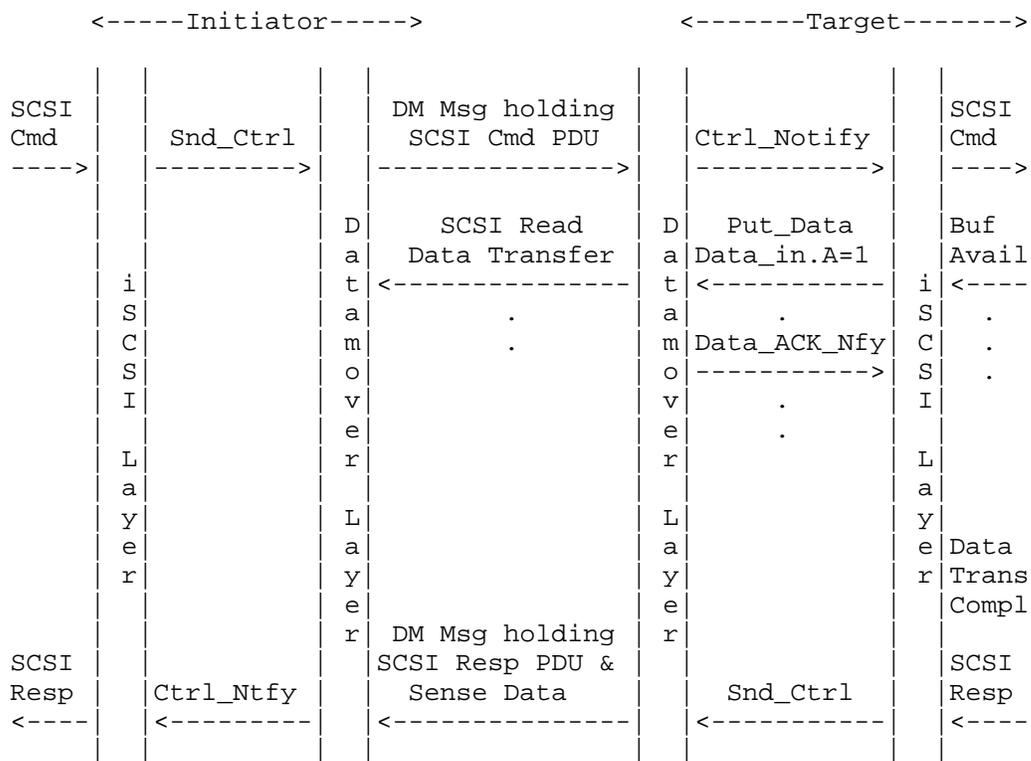


Figure 11. A SCSI Read Data Acknowledgement

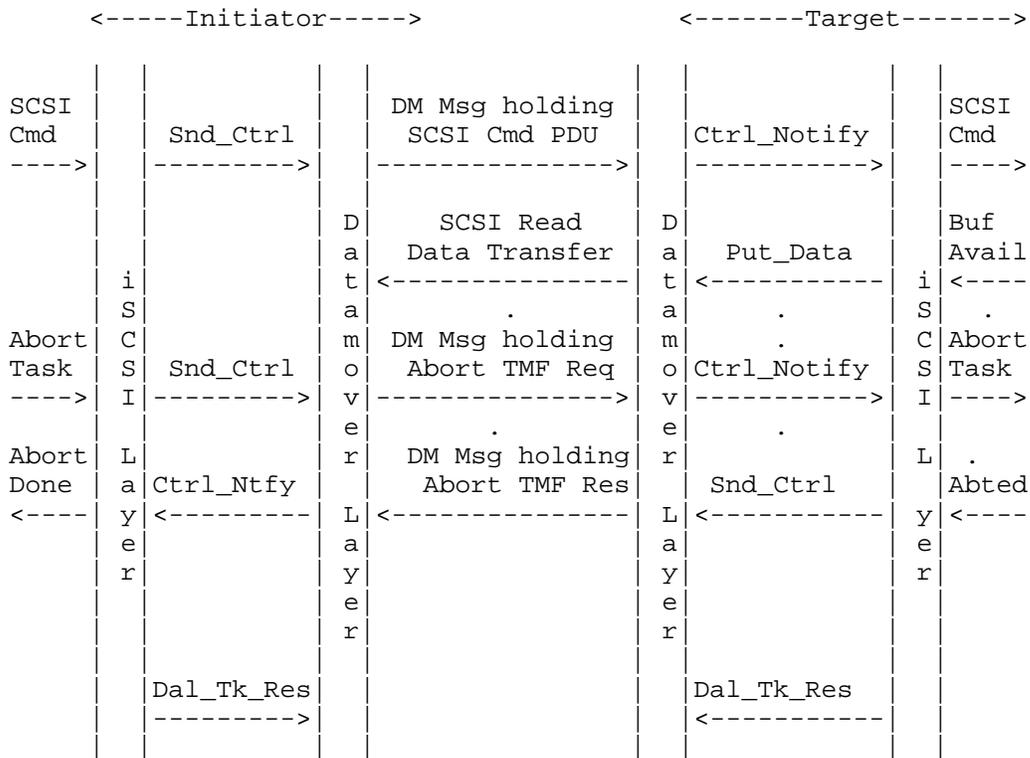


Figure 12. Task Resource Cleanup on Abort

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